

Corporate Technology

System and Application Analysis with LTTng

Project examples

- Serial input latency
- Sporadic delay in high prio application thread

How LTTng was used in Siemens projects to solve problems Gernot Hillier, CT SE 2

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Situation

- target hardware connected to serial port of "fast machine" (Multicore, fast 64bit CPUs, ...)
- high prio application has to react on signal from serial port within 10 ms
- expected time for reading character from serial port: << 1ms</p>

Problem

- •Unclear, undeterministic latency (several ms) from arrival of character to wakeup of application
- Possible reasons: locking, priority inversion, system load, ...?



Method:

- Trace simple testcase: bash shell prompt on serial port (does read(), write() in an endless loop to echo characters)
- Filter out events around IRQ arrival, then filter out right CPU

Trace output, part 1/2:

```
kernel_arch_syscall_entry: 308.506794 (cpu_3), 3051, bash, SYSCALL { syscall_id = 0 [sys_read+0x0/0xaa], ip = 0x7f29c1040f40 } => application blocks on read() [...]

kernel_irq_entry: 310.147709 (cpu_3), 0, swapper, IRQ { irq_id = 4, kernel_mode = 1, ip = 18446744071564210684 } => char arriving on serial line, IRQ 4 asserted kernel_timer_set: 310.147722 (cpu_3), 0, swapper, IRQ { expires = 4294944731, function = 0xfffffff80245288 [...] } => timer set, calling function delayed_work_timer_fn() kernel_irq_exit: 310.147728 (cpu_3), 0, swapper, SYSCALL { handled = 1 } => now CPU does nothing (no event) for nearly 3 ms
```



Trace output, part 2/2:

```
kernel_irq_entry: 310.150507 (cpu_3), 0, swapper, IRQ { irq_id = 239, kernel_mode = 1, ip =
18446744071564210684 } => timer interrupt occurs (HZ = 250)
[...]
kernel_sched_schedule: 310.150549 (cpu_3), 14, events/3, SYSCALL { prev_pid = 0,
next_pid = 14, prev_state = 0 } => workqueue thread for CPU 3 gets scheduled
kernel sched try wakeup: 310.150562 (cpu 3), 14, events/3, SYSCALL { pid = 3051, state
= 1 } => workqueue thread wakes up ...
kernel_sched_schedule: 310.150570 (cpu_3), 3051, bash, SYSCALL { prev_pid = 14,
next pid = 3051, prev state = 1 } => ... our target process
kernel arch syscall exit: 310.150585 (cpu 3), 3051, bash, USER MODE { ret = 1 }
=> read() returns, target application ...
kernel arch syscall entry: 310.150600 (cpu 3), 3051, bash, SYSCALL { syscall id = 1
[sys write+0x0/0xaa], ip = 0x7f29c1040fc0 } => ... can finally echo character
```



Further steps

- Code review of Linux serial driver, search for usage of work queues (serial8250 interrupt → serial8250 handle port → receive chars → tty flip buffer push → schedule delayed work → queue delayed work → queue delayed work on)
- Result: Linux serial code uses delayed work queues (delay: 1 jiffy) to handle incoming characters
- Reason: don't wake up userspace on each character => reduce overhead, increase throughput for "normal applications"

Conclusion

- throughput optimization in Linux serial code conflicts with our use case
- Own, simple serial receive routine was implemented for critical path



Situation

- multi-threaded application with a number of low-prio workers and highprio thread with "soft realtime" requirements
- large C++ application using rich middleware
- full code review (down to system call level) very time-consuming

Problem

- unclear, sporadic delays at some code points
- Possible reasons: system load, locking, application problems, ...?



Method:

- roughly identify problematic code points with application time stamping
- mark problematic code passages with invalid syscalls (syscall 600 at start of code point, 601 after code point if delay too high); LTTng userspace markers not ready
- Filter out events between syscalls, then filter out right CPU

Trace output, first try (excerpt):

```
kernel_arch_syscall_entry: 565.743 (cpu_5), 3309, PrioThread, { syscall_id = 600 }
kernel_arch_syscall_exit: 565.743 (cpu_5), 3309, PrioThread, { ret = -38 }
kernel sched schedule: 565.743 (cpu 5), 0, swapper, { prev pid = 3309, next pid = 0,
prev state = 2 } => kernel switches to idle task for unknown reason
kernel_arch_trap_entry: 565.774 (cpu_5), 3309, PrioThread, { trap_id = 14, ... } => process
returned from page fault after 30 ms!
kernel_arch_syscall_entry: 565.774 (cpu_5), 3309, PrioThread, { syscall_id = 601 }
```



Method:

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- Review kernel code => trace marker in page fault is too late
- Add additional trace markers to arch/x86/mm/fault.c:do_page_fault

Trace output, second try (excerpt):

```
kernel arch syscall entry: 575.226 (cpu 5), 3309, PrioThread, { syscall id = 600 }
kernel_arch_syscall_exit: 575.226 (cpu_5), 3309, PrioThread, { ret = -38 }
kernel_arch_page_fault_entry: 575.226 (cpu_5), 3309, PrioThread { ip = 0x7f1a3c073d9c }
=> theory confirmed, delay is caused by page fault
kernel arch page fault addr: 575.226 (cpu 5), 3309, PrioThread, { addr =
139735431535128, ip = 0x7f1a3c0 } => additional trace point to get faulting address
kernel_sched_schedule: 575.226 (cpu_5), 3336, WorkThread, { prev_pid = 3309, next_pid =
3336 } => This time, another low prio thread is runnable and scheduled
kernel_arch_trap_entry: 575.258 (cpu_5), 3309, PrioThread, { trap_id = 14, ... }
kernel_arch_syscall_entry: 575.258 (cpu_5), 3309, PrioThread, { syscall_id = 601 }
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```



Further steps

- decoding of page faulting address (with the help of /proc/<pid>/maps) => mmaped area on disk
- code review for access to mmaped disk area

Conclusion

- commonly used application function used by high prio thread caused access to mmaped area on disk
- application code restructuring to get rid of this access